

Math 551

2123

- HW #2 due Th 2/25
- HW #3 due Th 3/4
- Look over 3.1 - 3.4

# Floating Point Systems

For any  $\beta$   $b_1 \neq 0$

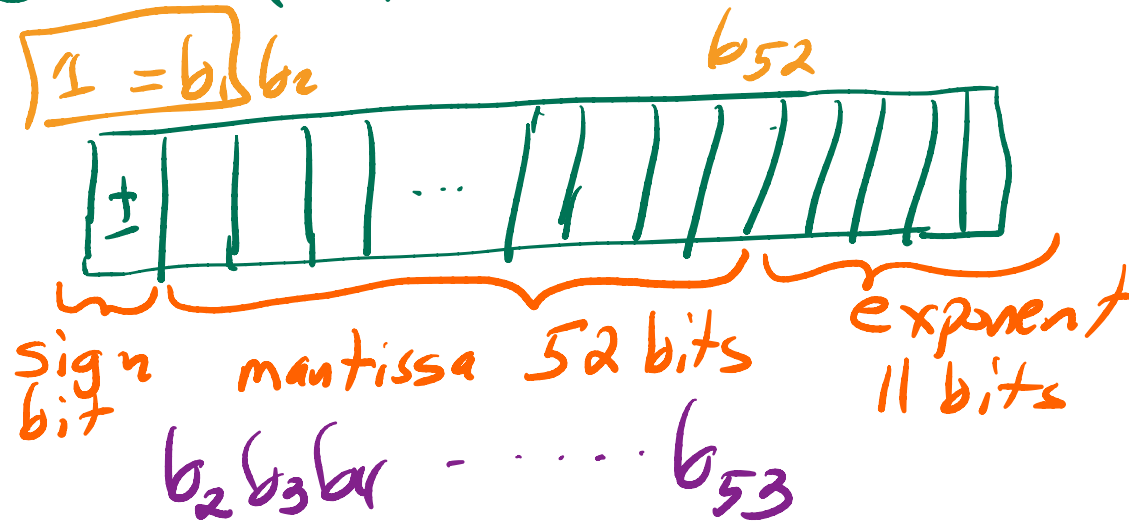
Double-Precision (64-bits IEEE)

$$\beta = 2$$

$$\{0, 1\}$$

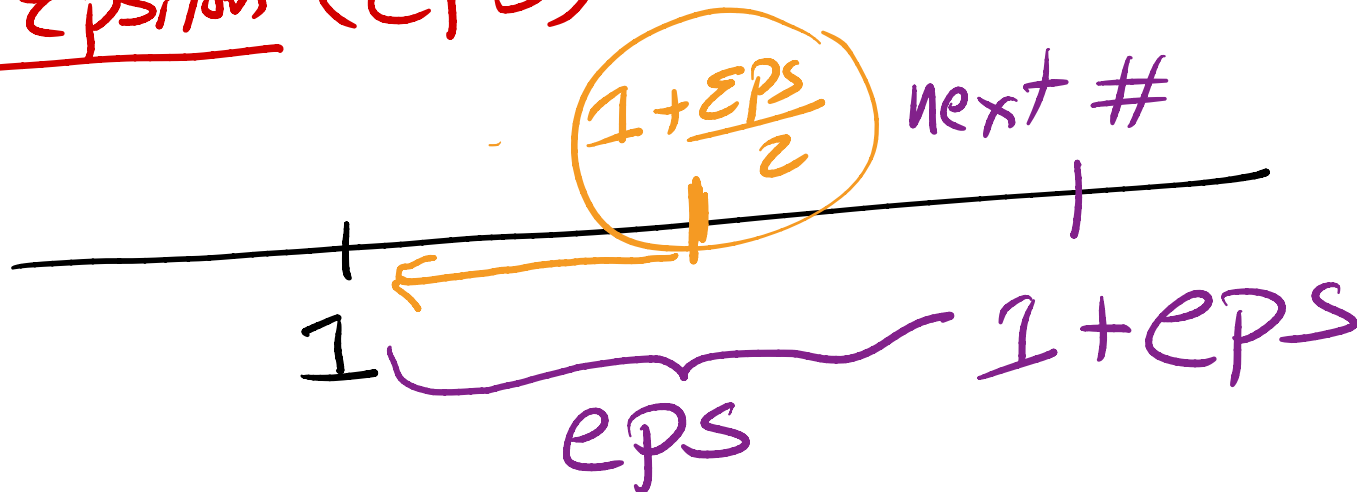
$$b_1 = 1$$

$$f(x) =$$



Machine Epsilon (eps)

NOT REPRESENTABLE!



eps?

$$f(1) = + (0.10000 \dots 00)_2 \times 2^1 = 1$$

$$f(1) + \text{eps} = + (0.\underbrace{100 \dots}_{b_1} \dots \underbrace{01}_{b_{52}})_2 \times 2^1$$

$$\text{so eps} = + (0.000 \dots \underbrace{001}_{2^{-52}})_2 \times 2^1$$

$$= (+)(2^{-52})(2) = \cancel{2^{-51}} \text{ Wrong!}$$

# MATLAB

NOT STORED

$$fl(1) + eps = + (0.\boxed{1}0 \dots 0 \underbrace{1}_{b_{53}})_2 \times 2^1$$

$$\Rightarrow eps = + (2^{-53}) \cdot 2 = + 2^{-52}$$

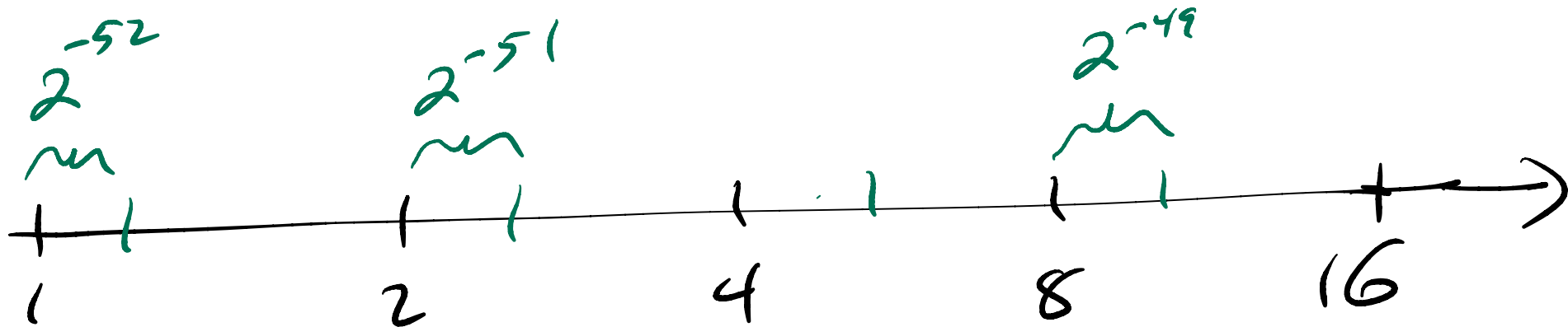
as we saw in MATLAB

$$F(8) + \text{eps}_8 = + \underbrace{(0.10 \dots 0001)_2}_{2^{-1} 2^{-2} \dots 2^{-52} 2^{-53}} \times \underbrace{2^4}_4$$

$$\text{eps}_8 = + (0.00 \dots 001)_2 \times 2^4$$

$$\underbrace{2^{-53}}_{2^{-53}} \cdot 2^4 = 2^{-49}$$

$$\underbrace{\quad \quad \quad}_{\text{eps}_8 = 2^{-49}} \quad 8 \quad \quad 8 + \underline{\text{eps}_8}$$



# Roots of Nonlinear Equations $F: \mathbb{R} \rightarrow \mathbb{R}$

Defn: Given  $f(x)$ ,  $\alpha$  is called a ROOT or ZERO of  $f(x)$  if  $f(\alpha) = 0$ .

Ex:  $F(x) = x^2 - 3x + 2 = (x-2)(x-1) = 0$   
 $\Rightarrow \alpha = 1, 2$

Ex:  $F(x) = 2 - e^x = 0$   
 $\Rightarrow e^x = 2$

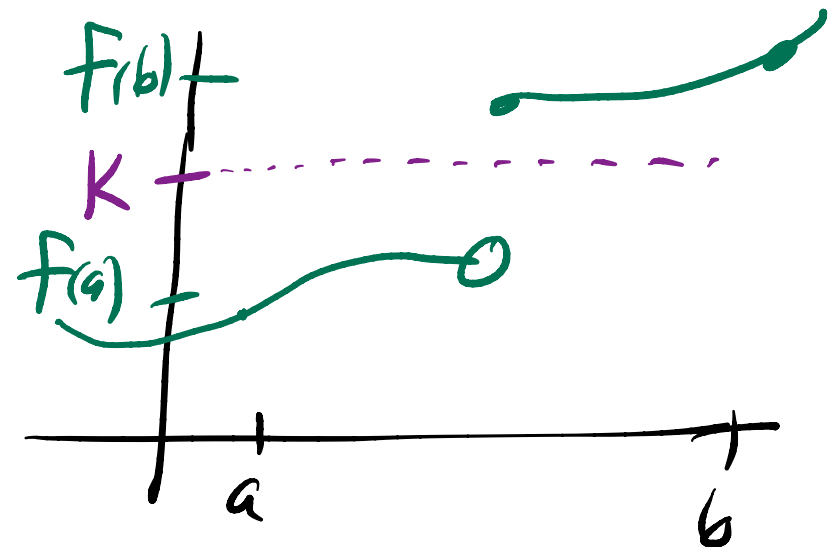
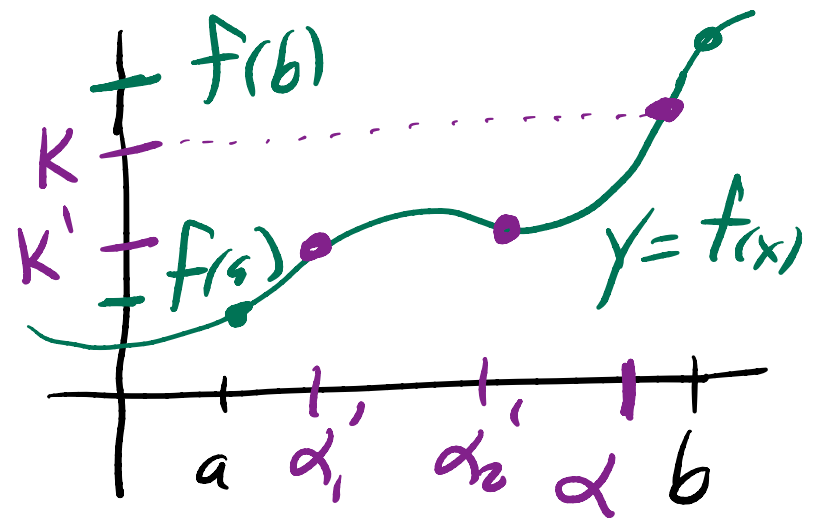
$\Rightarrow \alpha = \ln(2) = 0.69314 \dots$

Question : Given  $f(x)$  how do we determine the roots? Algorithm?

Thm (Intermediate Value Theorem)  $IVT$

Suppose  $f \in C([a, b])$  and  $K$  is any value between  $f(a)$  and  $f(b)$ . Then there exist  $\alpha \in (a, b)$  such that

$$f(\alpha) = K.$$

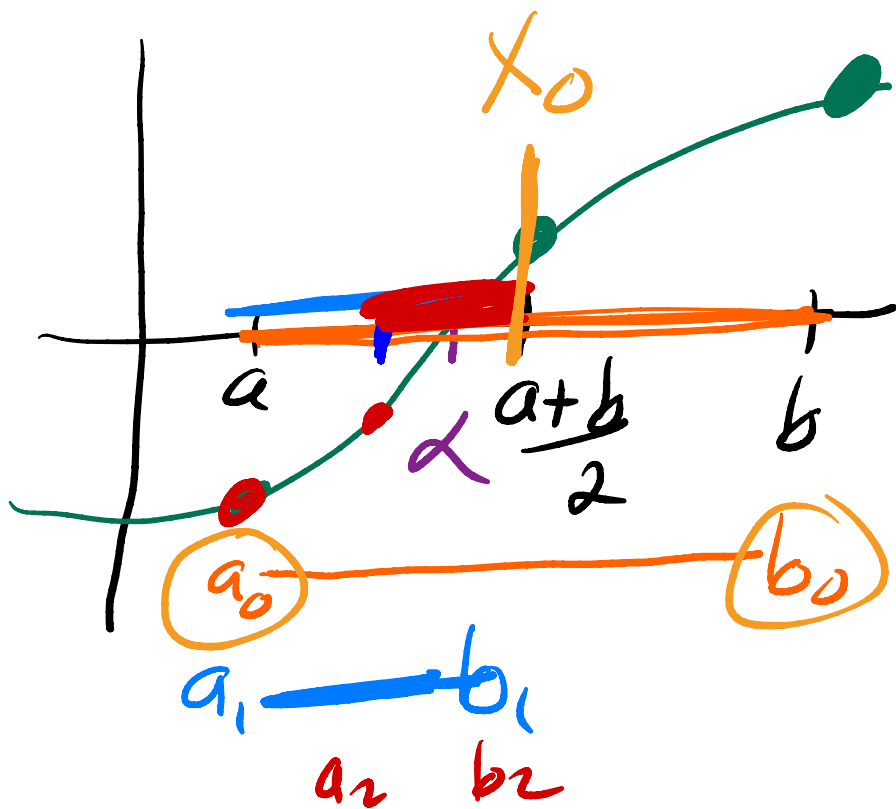




# Algorithm:

- ①  $f \in C([a, b])$
- ②  $f(a) \cdot f(b) < 0$   
 $f(a)$  and  $f(b)$  opposite sign

IUT  $\Rightarrow$  With  $K=0$  the  
IUT tells us there  
exists  $\alpha \in (a, b)$  such  
 $f(\alpha) = 0$ .



check the sign  
of  $f\left(\frac{a+b}{2}\right)$

$\alpha \in$   
 $(a_0, b_0)$   
 $(a_1, b_1)$   
 $(a_2, b_2)$

# Bisection

Given  $f \in C([a, b])$

$$f(a) \cdot f(b) < 0$$

Divide  
and Conquer  
Algorithm loop

$$a_0 = a$$

$$b_0 = b$$

$$\underline{n = 0}$$

$$x_n = (a_n + b_n) / 2$$

$$\text{if } f(a_n) \cdot f(x_n) < 0$$

$$a_{n+1} = a_n$$

$$b_{n+1} = x_n$$

Function  
eval

else

$$a_{n+1} = x_n$$

$$b_{n+1} = b_n$$

end

$$n = n + 1$$