1. (3.1-*Trefethen & Bau*) Prove that if W is an arbitrary nonsingular matrix, the function  $\|\cdot\|_W$  defined by  $\|x\|_W = \|Wx\|$  is a vector norm.

 $\underline{\mathbf{ANS}}$ : We need to show

(i)  $||x||_W \ge 0$ ,  $||x||_W = 0 \iff x = 0$ 

Given any  $x, ||x||_W = ||Wx|| \ge 0$  since  $||\cdot|| \ge 0$ . Let y = Wx, and note since W is nonsingular  $y = 0 \Leftrightarrow x = 0$ . Then  $||x||_W = ||y|| = 0 \Leftrightarrow y = 0 \Leftrightarrow x = 0$ .

- (ii)  $\|\alpha x\|_W = |\alpha| \|x\|_W$  $\|\alpha x\|_W = \|W(\alpha x)\| = \|\alpha W x\| = |\alpha| \|W x\|$ , since this is true for the  $\|\cdot\|$  norm. Thus,  $\|\alpha x\|_W = |\alpha| \|W x\| = |\alpha| \|x\|_W$ .
- (iii)  $||x + y||_W \le ||x||_W + ||y||_W$  $||x + y||_W = ||W(x + y)|| = ||Wx + Wy|| \le ||Wx|| + ||Wy||$ , since  $||u + v|| \le ||u|| + ||v||$  for any u, v. Thus,  $||x + y||_W \le ||Wx|| + ||Wy|| = ||x||_W + ||y||_W$ .

2. (3.2-Trefethen & Bau) Let  $\|\cdot\|$  denote any norm on  $\mathbb{C}^m$  and also the induced matrix norm on  $\mathbb{C}^{m \times m}$ . Show that  $\rho(A) \leq \|A\|$ , where  $\rho(A)$  is the spectral radius of A, i.e., the largest absolute value  $|\lambda|$  of an eigenvalue  $\lambda$  of A.

**<u>ANS</u>**: First note that for any x,  $||Ax|| \le ||A|| ||x||$  since this is a property of any induced matrix norm. Now, let  $(\lambda, x)$  be any e-pair of A, i.e.,  $Ax = \lambda x$ . Then

$$|\lambda|||x|| = ||\lambda x|| = ||Ax|| \le ||A|| ||x|| \quad \Rightarrow \quad |\lambda| \le ||A||.$$

Note we divided by  $||x|| \neq 0$  since x is and e-vector, so we must have  $x \neq 0$ . Thus  $|\lambda| \leq ||A||$ , and taking the sup over all  $\lambda$  on the left-handside (in fact a max since their are only a finite number of e-values), and noting the the right-handside is independent of  $\lambda$  gives the result.

- 3. (3.3-Trefethen & Bau) Vector and matrix p-norms are related by various inequalities, often involving the dimensions m or n. For each of the following, verify the inequality and give an example of a nonzero vector or matrix (for general m,n) for which equality is achieved. In this problem x is a m-vector and A is a  $m \times n$  matrix.
  - (a)  $||x||_{\infty} \le ||x||_{2}$ , (b)  $||x||_{2} \le \sqrt{m} ||x||_{\infty}$ ,
  - (c)  $||A||_{\infty} \le \sqrt{n} ||A||_{2}$ , (c)  $||A||_{\infty} \le \sqrt{n} ||A||_{2}$ ,
  - (d)  $||A||_2 < \sqrt{m} ||A||_\infty$ .

## ANS:

(a)  $||x||_{\infty} = \max_i |x_i| = \max_i (|x_i|^2)^{1/2} \le (\sum_{i=1}^m |x_i|^2)^{1/2} = ||x||_2$ . Equality is achieved for  $x = e_k, k = 1, \ldots, m$ , where  $e_k$  is one of the standard basis vectors for  $\mathbb{C}^m$ .

(b)  $\|x\|_2 = \left(\sum_{i=1}^m |x_i|^2\right)^{1/2} \le \left(\sum_{i=1}^m \max_i |x_i|^2\right)^{1/2} = \left(m\|x\|_{\infty}^2\right)^{1/2} = \sqrt{m}\|x\|_{\infty}$ . For equality take  $x = (1, 1, \dots, 1)^T \in \mathbb{C}^m$ . Then  $\|x\|_{\infty} = 1$  and  $\|x\|_2 = \sqrt{m}$ .

(c) We have

$$\begin{split} \|A\|_{\infty} &= \sup_{x \neq 0} \frac{\|Ax\|_{\infty}}{\|x\|_{\infty}} \\ &\leq \sup_{x \neq 0} \frac{\|Ax\|_2}{\|x\|_{\infty}} \quad \text{using (a) in the numerator} \\ &\leq \sup_{x \neq 0} \frac{\|Ax\|_2}{\|x\|_2/\sqrt{n}} \quad \text{using (b) in the denominator, and noting } x \in \mathbb{C}^n \\ &= \sqrt{n} \|A\|_2 \end{split}$$

For equality, let  $A \in \mathbb{C}^{m \times n}$  be the matrix whose first row is all ones, and zeros elsewhere. Clearly  $||A||_{\infty} = n$ . Now,  $A^*A$  is an  $n \times n$  matrix whose entries are all equal to one (check!) and its rank is 1. So 0 is an e-value of  $A^*A$  of multiplicity n-1. What is the remaining e-value? It is  $\lambda_n = n$ , which is easily seen to be the case since  $x = (1, 1, \ldots, 1)^T \in \mathbb{C}^n$  is a corresponding e-vector (check!). Thus  $||A||_2 = \sqrt{\rho(A^*A)} = \sqrt{n}$ . So we have  $||A||_{\infty} = n = \sqrt{n}\sqrt{n} = \sqrt{n}||A||_2$ . (d) We have

$$\begin{aligned} \|A\|_2 &= \sup_{x \neq 0} \frac{\|Ax\|_2}{\|x\|_2} \\ &\leq \sup_{x \neq 0} \frac{\sqrt{m} \|Ax\|_{\infty}}{\|x\|_2} \qquad \text{using (b) in the numerator, and noting } Ax \in \mathbb{C}^m \\ &\leq \sup_{x \neq 0} \frac{\sqrt{m} \|Ax\|_{\infty}}{\|x\|_{\infty}} \qquad \text{using (a) in the denominator} \\ &= \sqrt{m} \|A\|_{\infty} \end{aligned}$$

For equality, let  $A \in \mathbb{C}^{m \times n}$  be the matrix whose first column is all ones, and zeros elsewhere. Clearly  $||A||_{\infty} = 1$ . Now,  $A^*A$  is an  $n \times n$  diagonal matrix whose entries are all equal to zero, except for the (1, 1) entry, which is equal to m (check!). Thus,  $||A||_2 = \sqrt{\rho(A^*A)} = \sqrt{m}$ . So we have  $||A||_2 = \sqrt{m} = \sqrt{m} \times 1 = \sqrt{m} ||A||_{\infty}$ .

- 4. Prove that given a vector norm ||x||, the formula ||A|| = sup ||Ax|| / ||x|| defines a matrix norm for a square matrix A. Recall, this is referred to as the *induced matrix norm*.
  <u>ANS</u>: Suppose A ∈ C<sup>m×m</sup>. We need to show:
  - (i)  $||A|| \ge 0$ , and ||A|| = 0 only if A = 0.

 $||A|| \ge 0$  since it is the supremum of the ratio of  $||x||, ||Ax|| \ge 0$ . Now suppose ||A|| = 0but  $A \ne 0$ . Since  $A \ne 0$  there exists an  $x \ne 0$  such that  $Ax = y \ne 0$ . For example, if  $a_k \ne 0$  where  $a_k$  is the  $k^{th}$  column of A (and there must be one since  $A \ne 0$ ) let  $x = e_k$ . Then  $y = a_k$  and ||Ax||/||x|| = ||y||/||x|| > 0, hence so is the suprmum over all  $x \ne 0$ . Contradiction. So A = 0.

(ii)  $||A + B|| \le ||A|| + ||B||.$ 

We have  $||x + y|| \le ||x|| + ||y||$  for any  $x, y \in \mathbb{C}^m$ . Then

$$A + B \| = \sup_{x \neq 0} \frac{\|(A + B)x\|}{\|x\|} = \sup_{x \neq 0} \frac{\|Ax + Bx\|}{\|x\|}$$
$$\leq \sup_{x \neq 0} \frac{\|Ax\| + \|Bx\|}{\|x\|}$$
$$= \sup_{x \neq 0} \frac{\|Ax\|}{\|x\|} + \sup_{x \neq 0} \frac{\|Bx\|}{\|x\|}$$
$$= \|A\| + \|B\|.$$

(iii)  $\|\alpha A\| = |\alpha| \|A\|.$ 

We have  $\|\alpha x\| = |\alpha| \|x\|$  for any  $x \in \mathbb{C}^m$ ,  $\alpha \in \mathbb{C}$ . Then

$$\|\alpha A\| = \sup_{x \neq 0} \frac{\|\alpha Ax\|}{\|x\|} = \sup_{x \neq 0} \frac{|\alpha| \|Ax\|}{\|x\|}$$
$$= \sup_{x \neq 0} \frac{\alpha \|Ax\|}{\|x\|} = \alpha \sup_{x \neq 0} \frac{\|Ax\|}{\|x\|}$$
$$= |\alpha| \|A\|.$$

5. Prove that  $||A||_{\infty} = \max_{i} \sum_{j} |a_{ij}|$ , the maximum absolute row sum of the matrix A.

**<u>ANS</u>**: Assume  $A \in \mathbb{C}^{m \times n}$  and  $x \in \mathbb{C}^n$ . Then

$$||Ax||_{\infty} = \max_{i} |(Ax)_{i}| = \max_{i} |\sum_{j} a_{ij}x_{j}| = \max_{i} \sum_{j} |a_{ij}||x_{j}| \le ||x||_{\infty} \max_{i} \sum_{j} |a_{ij}|$$

So  $\frac{||Ax||_{\infty}}{||x||_{\infty}} \leq \max_{i} \sum_{j} |a_{ij}|$  for all  $x \neq 0$ . Now let k be such that  $\max_{i} \sum_{j} |a_{ij}| = \sum_{j} |a_{kj}|$ . If there is more than one such k choose the minimum. We can then write  $a_{kj} = r_j e^{i\theta_j}$  for some real number  $r_j \geq 0$  and  $0 \leq \theta_j < 2\pi$ . Note  $|a_{kj}| = |r_j e^{i\theta_j}| = |r_j| |e^{i\theta_j}| = |r_j| = r_j$ . Define  $\tilde{x} \in \mathbb{C}^n$  by  $\tilde{x}_j = e^{-i\theta_j}$  for  $j = 1, \ldots, n$ . Notice that is A were a real matrix then we would have  $\tilde{x}_j = \pm 1 = e^{-i(0,\pi)}$ . Then  $\|\tilde{x}\|_{\infty} = 1$  and

$$|\sum_{j} a_{kj} \tilde{x}_{j}| = |\sum_{j} r_{j} e^{i\theta_{j}} e^{-i\theta_{j}}| = |\sum_{j} r_{j}| = \sum_{j} r_{j} = \sum_{j} |a_{kj}|.$$

Finally,

$$||Ax||_{\infty} = \sup_{x \neq 0} \frac{||Ax||_{\infty}}{||x||_{\infty}} \le \max_{i} \sum_{j} |a_{ij}| = \sum_{j} |a_{kj}| = \frac{||A\tilde{x}||_{\infty}}{||\tilde{x}||_{\infty}} \le \sup_{x \neq 0} \frac{||Ax||_{\infty}}{||x||_{\infty}} = ||Ax||_{\infty}$$

Thus,  $||Ax||_{\infty} = \max_{i} \sum_{j} |a_{ij}|$ , the maximum absolute row sum of A.

## 6. Consider the 2-point BVP

$$\begin{cases} -y'' + (4x^2 + 2)y = 2x(1 + 2x^2) \\ y(0) = 1, \ y(1) = 1 + e \end{cases}$$

The exact solution is  $y(x) = x + e^{x^2}$ . Write a MATLAB function M-file to solve the problem using the **4th** order centered compact FD scheme

$$-D_{+}D_{-}\left(1-\frac{h^{2}}{12}c_{i}\right)u_{i}+c_{i}u_{i}=\left(1+\frac{h^{2}}{12}D_{+}D_{-}\right)f_{i}.$$

Use meshsize  $h = 1/2^p$ , where p is a positive integer. Your code should use your M-files **trilu** and **trilu\_solve**. For p = 1 : 4, plot the exact solution (y(x) vs. x) and the numerical solution  $(u_i \text{ vs. } x_i)$ , including the boundary points. The 4 plots should appear separately in one figure, with axes labeled and a title for each indicating p. Investigate **subplot** in MATLAB for how to have multiple plots in a single figure window. For p = 1 : 20 present a table with the following data - column 1: h; column 2:  $||u_h - y_h||_{\infty}$ ; column 3:  $||u_h - y_h||_{\infty}/h^4$ ; column 4: cpu time; column 5: (cpu time)/m, where h = 1/(m + 1). Discuss the trends in each column. Also, compare the accuracy for each h with the results of the 2nd order code from HW #1. How do the computational times compare? Include a copy of your code.

**<u>ANS</u>**: Writing out the discretization gives

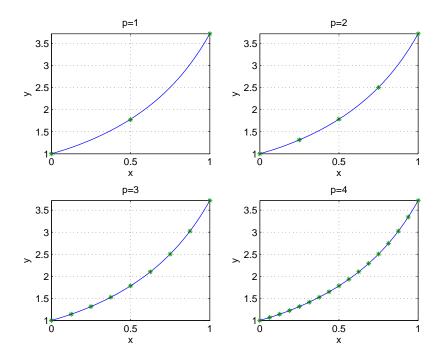
$$\frac{-(1-\frac{h^2}{12}c_{i-1})u_{i-1} + (2+\frac{5h^2}{6}c_i)u_i - (1-\frac{h^2}{12}c_{i+1})u_{i+1}}{h^2} = f_i + \frac{(f_{i-1}-2f_i+f_{i+1})}{12}$$

Here is the code for the first part of the problem, followed by a listing of the M-file function  $bvp\_solve4$ . The latter requires the M-files trilu and trilu\\_solve from problem 7 (hw #1).

```
xx=0:0.01:1;xx=xx';
yy=xx+exp(xx.^2);
c='4.*x.^2+2'; f='2*x.*(1+2*x.^2)';
clf;
for p=1:4
    [x,u]=bvp_solve4(2^p-1,0,1,1,1+exp(1),c,f);
    \max(abs(u-(x+exp(x.^2))))
    subplot(2,2,p),plot(xx,yy,x,u,'*'),grid
    axis('tight'),xlabel('x'),ylabel('y'),title(['p=',num2str(p)])
end
function [xv,uv,stime]=bvp_solve4(m,a,b,ya,yb,c,f)
%
h=(b-a)/(m+1);
xv=a:h:b; xv=xv';
fv=zeros(m+2,1);
cv=zeros(m+2,1);
```

```
%
x=xv;
cfv=eval(c);
fv=eval(f);
%
% setup tridiagonal matrix and rhs
%
av=(2+(5/6)*h^2*cfv(2:m+1))/(h^2);
bv=zeros(m,1); bv(2:m) =-(1-(h<sup>2</sup>/12)*cfv(2:m ))/(h<sup>2</sup>); bv(1)=0;
cv=zeros(m,1); cv(1:m-1)=-(1-(h<sup>2</sup>/12)*cfv(3:m+1))/(h<sup>2</sup>); cv(m)=0;
fv=fv(2:m+1)+(1/12)*(fv(1:m)-2*fv(2:m+1)+fv(3:m+2));
%
% BC adjustment
%
fv(1)=fv(1)+(1-(h^2/12)*cfv(1) )*ya/(h*h);
fv(m)=fv(m)+(1-(h^2/12)*cfv(m+2))*yb/(h*h);
%
% solve tridiagonal system; collect cpu time
%
tic;
[alpha,beta]=trilu(av,bv,cv);
uv=trilu_solve(alpha,beta,cv,fv);
stime=toc;
uv=[ya;uv;yb];
```

Here is the graph for p = 1, 2, 3 and 4.



Next we solve the BVP with  $m = 2^p - 1$  for p = 1, ..., 20. Here is the code:

```
c='4*x.^2+2'; f='2*x.*(1+2*x.^2)';
clf;
h=zeros(20,1); m=zeros(20,1);
times=zeros(20,1);
err_inf=zeros(20,1);
for p=1:20
    [x,u,stime]=bvp_solve4(2^p-1,0,1,1,1+exp(1),c,f);
    h(p)=1/(2^p);
    m(p)=2^p-1;
    times(p)=stime;
    y=x+exp(x.^2);
    err_inf(p)=max(abs(u-y));
end
format short e
disp(' ')
disp(' h inf_err err/h<sup>4</sup> cputime cputime/m ')
disp(' ------')
disp(' ')
disp([h err_inf err_inf./h.^4 times times./m])
```

The results are:

h	inf_err	err/h^4	cputime	cputime/m
5.0000e-01	1.2852e-02	2.0563e-01	9.2100e-04	9.2100e-04
2.5000e-01	9.7605e-04	2.4987e-01	1.9500e-04	6.5000e-05
1.2500e-01	6.3927e-05	2.6185e-01	1.3600e-04	1.9429e-05
6.2500e-02	4.0888e-06	2.6796e-01	1.3100e-04	8.7333e-06
3.1250e-02	2.5615e-07	2.6859e-01	1.5300e-04	4.9355e-06
1.5625e-02	1.6019e-08	2.6875e-01	2.1600e-04	3.4286e-06
7.8125e-03	1.0013e-09	2.6878e-01	2.5800e-03	2.0315e-05
3.9062e-03	6.1728e-11	2.6512e-01	6.3200e-04	2.4784e-06
1.9531e-03	1.0814e-12	7.4310e-02	1.0260e-03	2.0078e-06
9.7656e-04	1.3562e-11	1.4911e+01	2.2850e-03	2.2336e-06
4.8828e-04	5.5336e-11	9.7347e+02	4.4010e-03	2.1500e-06
2.4414e-04	2.1187e-10	5.9635e+04	8.8300e-03	2.1563e-06
1.2207e-04	8.7924e-10	3.9597e+06	1.6935e-02	2.0675e-06
6.1035e-05	1.2362e-09	8.9078e+07	3.5043e-02	2.1390e-06
3.0518e-05	2.3403e-09	2.6982e+09	6.8585e-02	2.0931e-06
1.5259e-05	1.3832e-09	2.5516e+10	1.3560e-01	2.0692e-06
7.6294e-06	1.3652e-08	4.0293e+12	2.6766e-01	2.0421e-06
3.8147e-06	4.3051e-07	2.0330e+15	5.3457e-01	2.0392e-06

1.9073e-06	1.3555e-07	1.0242e+16	1.0638e+00	2.0291e-06
9.5367e-07	7.0937e-07	8.5757e+17	2.1305e+00	2.0318e-06

We can see from the  $err/h^4$  column that the expected  $O(h^4)$  error is observed until  $h \approx 3.9062e - 3$  since  $err/h^4$  rapidly approaches a constant. But then we lose accuracy. Why? Roundoff error begins to dominate, as well as the conditioning of the matrix! Note, however, that the cputime stills scales linearly with m as evidenced by the cputime/m column.

Comparing with the second order method of HW #1 we see that the cpu timeS for the solvers are close. However, for moderate p the fourth order solution is clearly superior. For example, the error is 6.1728e - 1 for h = 3.9062e - 03, while for the second order method with the same h the error is 5.4521e - 06.

Note: I have only timed the linear solver portion of the code.