

## LINEAR RECURRENCES AND RATIONAL GENERATING FUNCTIONS

Let  $f(x) = \sum_k a_k x^k$  be a generating function. We say the sequence  $a_0, a_1, \dots$  satisfies a *linear recurrence relation of order  $n$*  if there exist constants  $c_1, \dots, c_n$  such that

$$(1) \quad a_k = c_1 a_{k-1} + \dots + c_n a_{k-n}$$

for all  $k \geq N$ , where  $N$  is some fixed number  $\geq n$ . The basic example is the Fibonacci sequence  $1, 1, 2, 3, 5, 8, 13, \dots$ , which is defined by  $a_0 = a_1 = 1$  and  $a_k = a_{k-1} + a_{k-2}$  for  $k \geq 2$ .

We have the following fact about the generating functions of sequences defined by recurrences:

**Theorem 1.** *Suppose the coefficients of  $f(x) = \sum a_k x^k$  satisfy (1) for all  $k \geq N$ , where  $N \geq n$  is given. Suppose further that the values  $a_i$  for  $i < N$  are known. Then*

$$(2) \quad f(x) = \frac{P(x)}{1 - c_1 x - c_2 x^2 - \dots - c_n x^n},$$

where  $P(x) = A_0 + A_1 x + \dots + A_{N-1} x^{N-1}$  is a polynomial of degree at most  $N - 1$ .

The coefficients of the numerator  $P(x)$  are determined by the  $a_i$ ,  $i < N$ , but in general  $A_i \neq a_i$ . A function of the form polynomial divided by polynomial is called a *rational function*. Sometimes people summarize the conclusion of the theorem by saying that the generating function of a sequence defined by a linear recurrence is a rational function.

Consider the Fibonacci example. According to the theorem the generating function for them can be written as

$$(3) \quad \frac{A_0 + A_1 x}{1 - x - x^2} = 1 + x + 2x^2 + 3x^3 + 5x^4 + 8x^5 + \dots$$

We can find the  $A_i$  by multiplying both sides of (3) by  $1 - x - x^2$ . There is a lot of cancellation (that's in fact the whole point of how this works), and we find

$$\begin{aligned} A_0 + A_1 x &= (1 - x - x^2)(1 + x + 2x^2 + 3x^3 + 5x^4 + 8x^5 + \dots) \\ &= 1 + 0x + 0x^2 + 0x^3 + \dots, \end{aligned}$$

which means  $A_0 = 1$  and  $A_1 = 0$ . Thus

$$\frac{1}{1 - x - x^2} = 1 + x + 2x^2 + 3x^3 + 5x^4 + 8x^5 + \dots$$

Here is how we can use an expression like (2) to get an explicit formula for the  $a_k$ . For simplicity we only discuss 2nd order recurrences with  $N = 2$  (the

general case is similar). Suppose one can break apart the rational function into “partial fractions”:

$$(4) \quad \frac{P(x)}{1 - c_1x - c_2x^2} = \frac{A}{1 - \alpha x} + \frac{B}{1 - \beta x}.$$

Here  $A, B$  are constants, and  $\alpha, \beta$  are the *inverses* of the roots of  $1 - c_1x - c_2x^2$ . (This is slightly different from the partial fractions one usually sees in calculus, because we want the denominators to look like  $1 - \alpha x$ , not  $x - \alpha \dots$  this reflects the fact that we use the inverses of the roots of the denominator of the left hand side.) Then one can conclude that

$$(5) \quad a_k = A\alpha^k + B\beta^k.$$

Why is (5) true? Because of the way the geometric series works:

$$(6) \quad \frac{1}{1 - \alpha x} = 1 + \alpha x + \alpha^2 x^2 + \alpha^3 x^3 + \dots$$

Plug (6) into (4) along with the analogous series for  $1/(1 - \beta x)$ , and then the coefficient of  $x^k$  is computed by the right hand side of (5).

For instance, consider the Fibonacci example. The roots of  $1 - x - x^2$  are  $r_1 = (\sqrt{5} - 1)/2$  and  $r_2 = (-\sqrt{5} - 1)/2$ . The inverses of these are  $1/r_2 = \alpha = (\sqrt{5} + 1)/2$  and  $1/r_1 = \beta = (-\sqrt{5} + 1)/2$ . So we get

$$(7) \quad a_k = A((\sqrt{5} + 1)/2)^k + B((-\sqrt{5} + 1)/2)^k.$$

We just have to solve for  $A$  and  $B$ . We can use the first two values of  $k$  for that:

$$\begin{aligned} 1 &= A + B \\ 1 &= A(\sqrt{5} + 1)/2 + B(-\sqrt{5} + 1)/2 \end{aligned}$$

It's somewhat painful to solve for  $A$  and  $B$ , but if we do we find

$$A = (5 + \sqrt{5})/10, \quad B = (5 - \sqrt{5})/10.$$

Here's an application of an expression like (7). Since  $|(-\sqrt{5} + 1)/2| < 1$  and  $(\sqrt{5} + 1)/2 > 1$ , it means that  $a_k$  must be very close to  $A((\sqrt{5} + 1)/2)^k$  for  $k$  large (because the other term in (7) will be close to 0). For example, when  $k = 20$ , the Fibonacci number is 10946, and

$$A((\sqrt{5} + 1)/2)^{100} = 10945.999981728 \dots$$

The relative error is  $10^{-9}$ . If  $k = 1000$  the relative error between the actual Fibonacci number and this approximation is  $10^{-419}$ , an unimaginably small number. How small is unimaginably small? The *Planck length*, which supposedly is the distance at which quantum effects begin to dominate in spacetime, is about  $10^{-35}$  m. The uncertainty principle of quantum mechanics means that we cannot measure distances smaller than this. Compared to our relative error, this number is gigantic.