

MATH 455 PROBLEM SET HINTS

PROBLEM SET 4

These are (usually) not complete solutions for the problems, but are intended to give you the basic ideas needed for a solution. If the basics of a problem are covered in class, either through working it out or doing a similar example, then we omit it here. Complete solutions typically involve more writing than is given here.

§1.4.1.

- (3) Only two land masses meet an odd number of bridges, so ...
- (4) Make a graph where the intersections are the vertices and the roads are the edges. Two vertices then have odd degree.

§1.4.2.

- (1) (a) An even cycle.
(b) Paste an odd and an even cycle together along a common vertex.
(c) Paste two odd cycles together along a common vertex.
(d) An odd cycle.
- (2) I'll leave the final answer to you.
- (6) $L(G)$ is connected if G is connected. $L(G)$ has a vertex for every edge and an edge for every pair of edges that meet in a vertex of G . If G is regular of degree r , then every vertex of $L(G)$ will meet $2r - 2$ edges in $L(G)$. Thus every vertex has even degree in $L(G)$.
- (7) I will leave (a) to you. For (b), we need every vertex to be even degree. This will happen if and only if n_1 and n_2 are even, since every vertex in K_{n_1, n_2} has degree n_1 or n_2 .

§1.5.1.

- (1) See Figures 1 and 2.
- (2) I did this by enlarging the polygon bounding a given region and "folding" the rest of the graph into the interior. Figure 3 shows the result for R_3 . (I added the vertex labels for my own benefit.)
- (7) See Figure 4 (sorry the labels are numbers, not letters).

§1.5.2.

- (1) Order 24 and regularity 3 implies that there are 36 edges. Since $v - e + f = 2$, we have 14 regions.
- (3) Any counterexample is good. For instance let G be two different vertices and no edges. Then $v - e + f = 2 - 0 + 1 = 3 \neq 2$.

- (10) One is given by the vertices and edges of the octahedron. But there are infinitely many 4-regular planar graphs. (If you want extra credit try to convince me of it with an original argument.)

§1.5.3.

- (1) $K_{2,2,2}$
- (2) Lots of possibilities. How about a soccer ball? Look on Wikipedia, say under *Archimedean solids*. In general you can take a bunch of random points on a sphere and then form their convex hull (the smallest convex set containing them) to make as many examples as you want.

§1.5.4.

- (2) This is challenging, but it can be done. Label the vertices of the outer pentagon as a, b, c, d, e starting from the 12 o'clock position and going clockwise, and label the vertices of the inner pentagon as A, B, C, D, E in the same fashion. Erase the edges bB and AC . The vertices A, B, C, b can now be erased to give $K_{3,3}$.
- (4) Assume that $n > 1$, since we know which complete graphs are planar. Then as soon as any two of the r_i are bigger than 2, the graph will contain $K_{3,3}$ as a subgraph and can't be planar. $K_{2,2}$ and $K_{2,2,2}$ are planar. What about $K_{2,2,2,2}$?

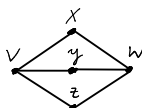


FIGURE 1.

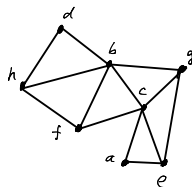


FIGURE 2.

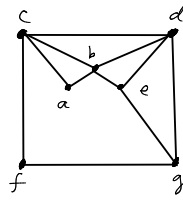


FIGURE 3.

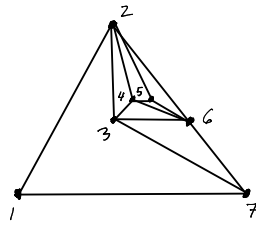


FIGURE 4.