

Tutte polynomial

(1)

For this topic, we allow loops and multiple edges in graphs.

The Tutte polynomial is a two variable poly attached to a graph G . Like the chromatic poly, it's defined via a "deletion-contraction" recursion.

Def The Tutte polynomial $T_G(x,y)$ is defined as follows:

- ① If G consists entirely of loops and bridges, then $T_G(x,y) = x^i y^j$, where i is the number of bridges and j is the number of loops.
- ② If e is an edge of G that is neither a loop nor a bridge, then

$$T_G(x,y) = T_{G-e}(x,y) + T_{G/e}(x,y)$$

Some examples

① $G = \text{Diagram of a cycle graph with 2 vertices} \Rightarrow T_G = xy^2$

② If G is a tree of order n , then $T_G = x^{n-1}$ (because all edges are bridges)

$$\textcircled{3} \quad G = K_3 = \begin{array}{c} \text{triangle} \end{array}$$

We illustrate the procedure graphically
 e ; not a bridge, not a loop

$$\begin{aligned} T\left(\begin{array}{c} \text{triangle} \end{array}\right) &= T(V) + T(O) \\ &= x^2 + T(\cdot) + T(\cdot) \\ &= x^2 + x + y. \end{aligned}$$

don't throw away multiple edge

$$\textcircled{4} \quad G = K_4 = \begin{array}{c} \text{tetrahedron} \end{array}$$

This takes some effort. The final answer is

$$x^3 + 3x^2 + 4xy + 2x + y^3 + 3y^2 + 2y.$$

We can do some preliminary computation first.

$$\textcircled{5} \quad T\left(\begin{array}{c} \text{loop} \end{array}\right) = x + y + y^2 + \dots + y^{n-1}$$

$\{\text{n-fold multiple edge}\}$

\textcircled{6} If G' is obtained from G by adding a loop, then $T_{G'} = y T_G$: $\textcircled{6}(G)$

\textcircled{7} Similarly, if G' is obtained from G by poking off an edge, then $T_{G'} = x T_G$:



Now we can try K_4 . (I'll omit the $T(\cdot)$ to simplify the notation, and will box graphs when we know their polynomials.) ③

$$\begin{aligned}
 A &= \text{Diagram of } K_4 = \text{Diagram with one edge removed} + \text{Diagram with one vertex removed} \\
 &= \text{Diagram with one edge removed} + \text{Diagram with two vertices removed} + \text{Diagram with three vertices removed} + \text{Diagram with four vertices removed} \\
 &= \text{Diagram with one edge removed} + 2(\text{Diagram with two vertices removed}) + \boxed{\text{Diagram with three vertices removed}} \\
 &= \boxed{\text{Diagram with one edge removed}} + \boxed{\text{Diagram with two vertices removed}} + 2(\boxed{\text{Diagram with two vertices removed}} + \boxed{\text{Diagram with three vertices removed}}) + \boxed{\text{Diagram with four vertices removed}}
 \end{aligned}$$

$$\textcircled{a} = x^3 \quad (\text{true})$$

$$\textcircled{b} = x^2 + x + y \quad (K_3)$$

$$2\textcircled{c} = 2xy + 2y^2 \quad (\text{loop on } \textcircled{d})$$

$$2\textcircled{d} = 2x^2 + 2xy \quad (\text{loop on } \textcircled{e})$$

$$\textcircled{e} = x + y + y^2 + y^3 \quad (4\text{-fold multiple edge})$$

$$x^3 + 3x^2 + 2x + 4xy + 2y + 3y^2 + y^3 \quad \checkmark$$

(4)

Connection between Tutte and the Chromatic polynomial

Let n be the order of G and $k(G)$ the number of connected components of G . Then

$$C_G(q) = (-1)^{n-k(G)} q^{k(G)} T_G(1-q, 0)$$

(we're using q instead of k for the variable)

Examples

① If G is a tree of order n , then $T_G(x,y) = x^{n-1}$.
 $k(G) = 1$, so we get

$$(-1)^{n-1} q^1 (1-q)^{n-1} = q(q-1)^{n-1} \checkmark$$

② If $G = K_3$, then $T_G = x^2 + x + y$. $n=3, k(G)=1$.
so we get

$$(-1)^2 q^1 ((q-1)^2 + (1-q) + 0)$$

$$= q^3 - 3q^2 + 2q = q(q-1)(q-2) \checkmark$$